

# CS450

## Structure of Higher Level Languages

Lecture 12: Finite-streams, evaluating expressions

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Press arrow keys   to change slides.

# Finite streams

# Finite streams

## The type `set`

A `set` is a finite stream of strings.

```
(define-type set
  ; A function that takes 0 args and builds either:
  (->
    ; an empty set OR an element and the rest of the stream set-add
    (U set-empty set-add)
  )
)
; set-empty denotes the end of the stream, so it has no fields
(struct set-empty ())
; set-add is akin to stream-add: holds a string and the rest of the stream
(struct set-add ([first : String] [rest : set]))
```

# An example of a finite stream

Here is an example of a set  $\{ "a", "b", "c" \}$

```
(thunk
  (set-add "a"
    (thunk
      (set-add "b"
        (thunk
          (set-add "c" (thunk (set-empty))))))))))
```

which is similar to building a list, but with `thunks` interleaved

```
(cons "a"
  (cons "b"
    (cons "c" empty)))
```

# Printing the elements of a finite-stream

```
(: print-set (-> set Void))
```

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```
(: print-set (-> set Void))

(define (print-set s)
  (match (s)
    ; Call s to build the next element:
    [(set-empty) (void)] ; If we there are no more elements, void does nothing
    [(set-add h l) ; Otherwise,
     (displayln h) ; print the element
     (print-set l) ; and loop
    ]
  )
)
```

# Copying a finite stream

```
(: copy (-> set set))
```

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```

## Finite streams

```
(define (copy l)
  (thunk
    (match (l)
      [(set-empty) (set-empty)]
      [(set-add h l)
       (set-add h (copy l))]
      ]
    )
  )
)
```

## Lists

```
(define (copy l)
  (match l
    [(list) (list)]
    [(list h l ...)
     (cons h (copy l))]
    ]
  )
)
```

Similarly, to infinite streams, when building a stream from another, we must **thunk** BEFORE we **match**.





# Faulty copy implementation

```
(: faulty-copy-set (-> set set))
(define (faulty-copy-set s)
  (match (s)
    [(set-empty) (thunk (set-empty))]
    [(set-add h s)
     (thunk (set-add h (faulty-copy-set s)))]))
```

- Calling `(faulty-copy-set s)` asks `s` to build the first element *before* building its copy.
- The problem is that when you ask `(faulty-copy-set s)` to build its *first* element, you are actually requesting `s` to build its *second* element.
- **Solution:** a `(thunk ...)` **must** surround any `(match (s) ...)`

# Understanding stream-fold

# Understanding stream-fold

Stream-fold behaves like the following function, which returns a list of the updates to the accumulator. Use this function as the blueprint to your solution.

```
(define (fold2 f a l)
  (match l
    [(list) (list)]
    [(list h l ...)
     ; h = e1
     ; l = [e2, e3]
     ; a1 = f e1 a
     (define a1 (f h a))
     (displayln a1)
     (define result (fold2 f a1 l))
     ; result = [a1, a2, a3]
     (cons a result)
     ; expected output: [a, f e1 a, f e2 (f e1 a), f e3 (f e2 (f e1 a))]
    ]))
```

# Evaluating expressions

# Evaluating expressions

Our goal is to implement an evaluation function that takes an expression and yields a value.

```
expression = value | variable | function-call  
value = number  
function-call = ( expression+ )
```

# How do we evaluate an expression

What is an expression?

*expression = value | variable | function-call*

How do we evaluate a value?

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```
expression = value | variable | function-call
```

How do we evaluate a value? **The evaluation of a value  $v$  is  $v$  itself.**

```
(check-equal? 10 (eval-exp (r:number 10)))
```

How do we evaluate a function call?

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How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**



# Example

How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**

```
(eval-exp  
  '(-  
    (+ 3 2)  
    (* 5 2)) )
```

```
①  
<- evaluate '-  
<- evaluate '(+ 3 2)  
<- evaluate '(* 5 2)
```

# Example

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```
(eval-exp  
  '(-  
    (+ 3 2)  
    (* 5 2)) )
```

```
= ((eval-exp '-)  
   (eval-exp '(+ 3 2))  
   (eval-exp '(* 5 2)))
```

```
①  
<- evaluate '-  
<- evaluate '(+ 3 2)  
<- evaluate '(* 5 2)
```

```
②  
<- evaluate '+, evaluate 3, evaluate 2  
<- evaluate '*, evaluate 5, evaluate 2
```

# Example

How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**

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(eval-exp  
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```

```
= ((eval-exp '-)  
   (eval-exp '(+ 3 2))  
   (eval-exp '(* 5 2)))
```

```
= ((eval-exp '-)  
   ((eval-exp '+) 3 2)  
   ((eval-exp '*) 5 2))
```

```
①  
<- evaluate '-  
<- evaluate '(+ 3 2)  
<- evaluate '(* 5 2)
```

```
②  
<- evaluate '+, evaluate 3, evaluate 2  
<- evaluate '*, evaluate 5, evaluate 2
```

```
③  
<- numbers are values, so just return those  
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```

# How do we evaluate arithmetic operators?

```
= ((eval-exp '-)  
   ((eval-exp '+) 3 2)  
   ((eval-exp '*) 5 2))
```

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```
= ((eval-exp '-)  
  ((eval-exp '+) 3 2)  
  ((eval-exp '*) 5 2))
```

```
= (-  
  (+ 3 2)  
  (* 5 2))
```

```
<- Evaluate '-' as function -  
<- Evaluate '+' as function +  
<- Evaluate '*' as function *
```

# Evaluation of arithmetic expressions

1. When evaluating a number, just return that number
2. When evaluating an arithmetic symbol, return the respective arithmetic function
3. When evaluating a function call evaluate each expression and apply the first expression to remaining ones

Essentially evaluating an expression **translates** our AST nodes as a Racket expression.

Implementing eval-exp...

# Specifying eval-exp

- We are use the AST we defined in Lesson 5, not datums.
- Assume function calls are binary.

```
(check-equal? (r:eval-exp (r:number 5)) 5)
(check-equal? (r:eval-exp (r:number 10)) 10)
(check-equal? (r:eval-exp (r:variable? '+)) +)
(check-equal?
  (r:eval-exp
    (r:apply
      (r:variable '+)
      (list (r:number 10) (r:number 5))))
  15)
```



# Implementing `eval-exp`

We are using the AST as structs, not datums. Assume function calls are binary.

```
(: r:eval-exp (-> r:expression Number))
(define (r:eval-exp exp)
  (match exp
    ; If it's a number, return that number
    [(r:number v) v]
    ; If it's a function with 2 arguments
    [(r:apply (r:variable f) (list arg1 arg2))
     (define func (r:eval-builtin f))
     (func (r:eval-exp arg1) (r:eval-exp arg2))]
  )
)
```

# Implementing `r:eval-builtin`

## Spec

```
(check-equal? (r:eval-builtin '+) +)  
(check-equal? (r:eval-builtin '-') -)
```

# Implementing `r:eval-builtin`

## Spec

```
(check-equal? (r:eval-builtin '+) +)
(check-equal? (r:eval-builtin '-') -)
```

## Solution

```
(define (r:eval-builtin sym)
  (cond [(equal? sym '+) +]
        [(equal? sym '*) *]))
```

# Handling functions with an arbitrary number of parameters

(required for Homework 4)

# Function `apply`

Function `(apply f args)` applies function `f` to the list of arguments `args`.

## Examples

```
(check-equal? (apply + (list 1 2 3 4)) 10)
```

Example: implement `(sum l)` that takes returns the summation of all members in `l` using `apply`.

## Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

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Function `(apply f args)` applies function `f` to the list of arguments `args`.

## Examples

```
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Example: implement `(sum l)` that takes returns the summation of all members in `l` using `apply`.

## Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

## Solution

```
(define (sum l) (apply + l))
```

# Handling multiple-args without apply

Some multi-arg operations can be implemented without the need of `apply`.

Implement `(sum 1)` without using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

# Handling multiple-args without apply

Some multi-arg operations can be implemented without the need of `apply`.

Implement `(sum 1)` without using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

Solution

```
(define (sum 1)
  (cond
    [(empty? 1) 0]
    [else (+ (first 1) (sum (rest 1)))]))
```



# Implementing functions with multi-args

How could we implement a function with multiple parameters, similar to `+`? **Use the `.` notation.**

The dot `.` notation declares that the next variable represents a list of zero or more parameters.

## Examples

```
(define (map-ex f . args)
  (map f args))

(check-equal? (list 2 3 4) (map-ex (curry + 1) 1 2 3))

(define (sum . 1) (foldl + 0 1))
(check-equal? 6 (sum 1 2 3))
```