

CS450

## Structure of Higher Level Languages

Lecture 10: Map, zip, enumerate, filter, expression evaluation

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# The map stream

# Map for streams

| Given a stream `s` defined as

`e0 e1 e2 e3 e4 ...`

| and a function `f` the stream `(stream-map f s)` should yield

`(f e0) (f e1) (f e2) (f e3) (f e4) ...`

# Map for streams

## Spec

```
#lang racket
(require rackunit)

(define s0
  (stream-map (curry + 2) (naturals)))
(check-equal? (stream-get s0) 2)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 3)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```



# Map for streams

## Spec

```
#lang racket
(require rackunit)

(define s0
  (stream-map (curry + 2) (naturals)))
(check-equal? (stream-get s0) 2)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 3)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```

## Solution

```
(define (stream-map f s)
  (define (stream-map-iter s)
    (cons
      (f (stream-get s))
      (thunk (stream-map-iter (stream-next s)))))

  (stream-map-iter s))
```



# The stream of even numbers

# Even naturals

Build a stream of even numbers. Tip: use `stream-map` and `naturals`.

0 2 4 6 8 10 12 ...

Spec

```
#lang racket
(require rackunit)
(define s0 (even-naturals))
(check-equal? (stream-get s0) 0)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 2)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```

# Even naturals

Build a stream of even numbers. Tip: use `stream-map` and `naturals`.

0 2 4 6 8 10 12 ...

Spec

```
#lang racket
(require rackunit)
(define s0 (even-naturals))
(check-equal? (stream-get s0) 0)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 2)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```

Solution

```
(define (even-naturals)
  (stream-map
    (curry * 2)
    (naturals)))
```



# Merge two streams

# Zip two streams

| Given a stream `s1` defined as

`e1 e2 e3 e4 ...`

| and a stream `s2` defined as

`f1 f2 f3 f4 ...`

| the stream (`stream-zip s1 s2`) returns

`(cons e1 f1) (cons e2 f2) (cons e3 f3) (cons e4 f4) ...`

# Zip for streams

## Spec

```
#lang racket
(require rackunit)
(define s0
  (stream-zip (naturals) (even-naturals)))

(check-equal? (stream-get s0) (cons 0 0))

(define s1 (stream-next s0))
(check-equal? (stream-get s1) (cons 1 2))

(define s2 (stream-next s1))
(check-equal? (stream-get s2) (cons 2 4))
```



# Zip for streams

## Spec

```
#lang racket
(require rackunit)
(define s0
  (stream-zip (naturals) (even-naturals)))
(check-equal? (stream-get s0) (cons 0 0))
(define s1 (stream-next s0))
(check-equal? (stream-get s1) (cons 1 2))
(define s2 (stream-next s1))
(check-equal? (stream-get s2) (cons 2 4))
```

## Solution

```
(define (stream-zip s1 s2)
  (define (stream-zip-iter s1 s2)
    (cons
      (cons (stream-get s1)
            (stream-get s2))
      (thunk
        (stream-zip-iter
          (stream-next s1)
          (stream-next s2))))))
(stream-zip-iter s1 s2))
```



# Exercises on streams

# Zip two streams

| Given a stream `s1` defined as

`e1 e2 e3 e4 ...`

| and a stream `s2` defined as

`f1 f2 f3 f4 ...`

| the stream (`stream-zip s1 s2`) returns

`(cons e1 f1) (cons e2 f2) (cons e3 f3) (cons e4 f4) ...`

# Enumerate a stream

| Build a stream from a given stream `s` defined as

`e0 e1 e2 e3 e4 e5 ...`

| the stream (`stream-enum s`) returns

`(cons 0 e0) (cons 1 e1) (cons 2 e2) (cons 3 e3) (cons 4 e4) (cons 5 e5) ...`

# Enumerate a stream

## Spec

```
#lang racket
(require rackunit)

(define s0 (stream-enumerate (even-naturals)))
(check-equal? (stream-get s0) (cons 0 0))

(define s1 (stream-next s0))
(check-equal? (stream-get s1) (cons 1 2))

(define s2 (stream-next s1))
(check-equal? (stream-get s2) (cons 2 4))
```



# Enumerate a stream

## Spec

```
#lang racket
(require rackunit)

(define s0 (stream-enum (even-naturals)))
(check-equal? (stream-get s0) (cons 0 0))

(define s1 (stream-next s0))
(check-equal? (stream-get s1) (cons 1 2))

(define s2 (stream-next s1))
(check-equal? (stream-get s2) (cons 2 4))
```

## Solution

```
(define (stream-enum s)
  (stream-zip (naturals) s))
```



# Filter

How would a filter work with streams?

# Filter

## Spec

```
#lang racket
(define s0
  (stream-filter (curry ≤ 10)
    (naturals)))
(check-equal? (stream-get s0) 10)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 11)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 12)
```



# Converting filter to stream-filter

```
; List version -----
1 (define (filter to-keep? l)
2  (cond
3   [(empty? l) l]
4   [(to-keep? (first l))
5    (cons (first l)
6          (filter to-keep? (rest l)))]
8   [else (filter to-keep? (rest l))]))
; Stream-version -----
1 (define (stream-filter to-keep? s)
2  (cond
3   ; ← no base case; streams are infinite
4   [(to-keep? (stream-get s)) ; ← first becomes stream-get
5    (cons (stream-get s)
6          ; Second element is always a thunk
7          (thunk (stream-filter to-keep? (stream-next s))))]
8   [else (stream-filter to-keep? (stream-next s))])) ; rest becomes stream-next
```



# Drop every other element

Given a stream defined below, drop every other element from the stream. That is, given a stream `s` defined as...

`e0 e1 e2 e3 e4 ...`

stream (`stream-drop-1 s`) returns

`e0 e2 e4 ...`

# Drop every other element...

Spec

```
#lang racket
(require rackunit)

(define s0 (stream-drop-1 (naturals)))
(check-equal? (stream-get s0) 0)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 2)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```



# Drop every other element...

## Spec

```
#lang racket
(require rackunit)

(define s0 (stream-drop-1 (naturals)))
(check-equal? (stream-get s0) 0)

(define s1 (stream-next s0))
(check-equal? (stream-get s1) 2)

(define s2 (stream-next s1))
(check-equal? (stream-get s2) 4)
```

## Solution

```
(define (stream-drop-1 s)
  ; for each e yield (i, e)
  (define enum-s (stream-enum s))
  ; given (i, e) only keep (even? i)
  (define even-s
    (stream-filter
      ;(lambda (x) (even? (car x)))
      (compose even? car)
      enum-s))
  ; convert (i, e) back to e
  (stream-map cdr even-s))
```



# More exercises

- `(stream-ref s n)` returns the element in the `n`-th position of stream `s`
- `(stream-interleave s1 s2)` interleave each element of stream `s1` with each element of `s2`
- `(stream-merge f s1 s2)` for each `i`-th element of stream `s1` (say `e1`) and `i`-th element of stream `s2` (say `e2`) return `(f e1 e2)`
- `(stream-drop n s)` ignore the first `n` elements from stream `s`
- `(stream-take n s)` returns the first `n` elements of stream `s` in a list in appearance order

# Evaluating expressions

# Evaluating expressions

Our goal is to implement an evaluation function that takes an expression and yields a value.

```
expression = value | variable | function-call  
value = number  
function-call = ( expression+ )
```



# How do we evaluate an expression

What is an expression?

```
expression = value | variable | function-call
```

| How do we evaluate a value?

# How do we evaluate an expression

What is an expression?

```
expression = value | variable | function-call
```

| How do we evaluate a value? **The evaluation of a value v is v itself.**

```
(check-equal? 10 (eval-exp (r:number 10)))
```

| How do we evaluate a function call?

# How do we evaluate an expression

What is an expression?

```
expression = value | variable | function-call
```

| How do we evaluate a value? **The evaluation of a value v is v itself.**

```
(check-equal? 10 (eval-exp (r:number 10)))
```

| How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**



# Example

How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**

```
(eval-exp
  '(-
    (+ 3 2)
    (* 5 2))) )
```

①  
← evaluate '-  
← evaluate '(+ 3 2)  
← evaluate '(\* 5 2)

# Example

How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**

```
(eval-exp
  '(-
    (+ 3 2)
    (* 5 2)))
```

```
= ((eval-exp '-')
  (eval-exp '(+ 3 2))
  (eval-exp '(* 5 2)))
```

①  
← evaluate '  
← evaluate '(+ 3 2)  
← evaluate '(\* 5 2)

②  
← evaluate '+, evaluate 3, evaluate 2  
← evaluate '\*', evaluate 5, evaluate 2

# Example

How do we evaluate a function call? **The evaluation of a function call evaluates each expression from left to right and then it applies the function to the arguments.**

```
(eval-exp
  '(-
    (+ 3 2)
    (* 5 2)) ))  
  
= ((eval-exp '-)
    (eval-exp '(+ 3 2))
    (eval-exp '(* 5 2)))  
  
= ((eval-exp '-)
    ((eval-exp '+) 3 2)
    ((eval-exp '*) 5 2))
```

- ①  
  ← evaluate '-'  
  ← evaluate '(+ 3 2)  
  ← evaluate '(\* 5 2)
  
- ②  
  ← evaluate '+, evaluate 3, evaluate 2  
  ← evaluate '\*', evaluate 5, evaluate 2
  
- ③  
  ← numbers are values, so just return those  
  ← numbers are values, so just return those

# How do we evaluate arithmetic operators?

```
= ((eval-exp '-)  
  ((eval-exp '+) 3 2)  
  ((eval-exp '*) 5 2))
```



# How do we evaluate arithmetic operators?

```
= ((eval-exp '-)  
  ((eval-exp '+) 3 2)  
  ((eval-exp '*)) 5 2))
```

```
= (-  
  (+ 3 2)  
  (* 5 2))
```

← Evaluate '-' as function -  
← Evaluate '+' as function +  
← Evaluate '\*' as function \*



# Evaluation of arithmetic expressions

1. When evaluating a number, just return that number
2. When evaluating an arithmetic symbol, return the respective arithmetic function
3. When evaluating a function call evaluate each expression and apply the first expression to remaining ones

Essentially evaluating an expression **translates** our AST nodes as a Racket expression.



# Implementing eval-exp...

# Specifying eval-exp

- We are use the AST we defined in Lesson 5, not datums.
- Assume function calls are binary.

```
(check-equal? (r:eval-exp (r:number 5)) 5)
(check-equal? (r:eval-exp (r:number 10)) 10)
(check-equal? (r:eval-exp (r:variable? '+)) +)
(check-equal?
  (r:eval-exp
    (r:apply
      (r:variable '+)
      (list (r:number 10) (r:number 5))))
```

15)



# Implementing eval-exp

We are using the AST we defined in Lesson 5, not datums. Assume function calls are binary.

```
(define (r:eval-exp exp)
  (cond
    ; 1. When evaluating a number, just return that number
    [(r:number? exp) (r:number-value exp)]
    ; 2. When evaluating an arithmetic symbol,
    ;     return the respective arithmetic function
    [(r:variable? exp) (r:eval-builtin (r:variable-name exp))]
    ; 3. When evaluating a function call evaluate each expression and apply
    ;     the first expression to remaining ones
    [(r:apply? exp)
      ((r:eval-exp (r:apply-func exp))
       (r:eval-exp (first (r:apply-args exp)))
       (r:eval-exp (second (r:apply-args exp))))]
    [else (error "Unknown expression:" exp)]))
```

# Implementing r:eval-builtins

## Spec

```
(check-equal? (r:eval-builtin '+) +)
(check-equal? (r:eval-builtin '-) -)
(check-equal? (r:eval-builtin '/) /)
(check-equal? (r:eval-builtin '*) *)
(check-equal? (r:eval-builtin 'foo) #f)
```



# Implementing r:eval-builtin

## Spec

```
(check-equal? (r:eval-builtin '+) +)
(check-equal? (r:eval-builtin '-) -)
(check-equal? (r:eval-builtin '/') /)
(check-equal? (r:eval-builtin '*)) *)
(check-equal? (r:eval-builtin 'foo) #f)
```

## Solution

```
(define (r:eval-builtin sym)
  (cond [(equal? sym '+) +]
        [(equal? sym '*)) *)
        [(equal? sym '-) -]
        [(equal? sym '/') /]
        [else #f]))
```



# Handling functions with an arbitrary number of parameters (required for Homework 3)

# Function apply

Function `(apply f args)` applies function `f` to the list of arguments `args`.

Examples

```
(check-equal? (apply + (list 1 2 3 4)) 10)
```

Example: implement `(sum l)` that takes returns the summation of all members in `l` using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```



# Function apply

Function `(apply f args)` applies function `f` to the list of arguments `args`.

Examples

```
(check-equal? (apply + (list 1 2 3 4)) 10)
```

Example: implement `(sum l)` that takes returns the summation of all members in `l` using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

Solution

```
(define (sum l) (apply + l))
```



# Handling multiple-args without apply

| Some multi-arg operations can be implemented without the need of `apply`.

Implement `(sum 1)` without using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

# Handling multiple-args without apply

| Some multi-arg operations can be implemented without the need of `apply`.

Implement `(sum 1)` without using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

Solution 1

```
(define (sum 1)
  (cond
    [(empty? 1) 0]
    [else (+ (first 1) (sum (rest 1))))]))
```

Solution 2 (foldl is tail-recursive)



# Handling multiple-args without apply

| Some multi-arg operations can be implemented without the need of `apply`.

Implement `(sum 1)` without using `apply`.

Spec

```
(check-equal? (sum (list)) 0)
(check-equal? (sum (list 1 2 3 4)) 10)
```

Solution 1

```
(define (sum 1)
  (cond
    [(empty? 1) 0]
    [else (+ (first 1) (sum (rest 1))))]))
```

Solution 2 (`foldl` is tail-recursive)

```
(define (sum 1) (foldl + 0 1))
```



# Implementing functions with multi-args

How could we implement a function with multiple parameters, similar to `+`? Use the `.` notation.

The dot `.` notation declares that the next variable represents a list of zero or more parameters.

## Examples

```
(define (map-ex f . args)
  (map f args))

(check-equal? (list 2 3 4) (map-ex (curry + 1) 1 2 3))
```

```
(define (sum . l) (foldl + 0 l))
(check-equal? 6 (sum 1 2 3))
```

