CS450

Structure of Higher Level Languages

Lecture 1: Course info, arithmetic in Racket, evaluation

Tiago Cogumbreiro

About the course

- **Instructor:** Tiago (蒂亚戈) Cogumbreiro (he/him)
- **Location:** (Y02-2300) Room 2300, 2nd floor, University Hall
- **Schedule:** Monday, Wednesday / 5:30PM - 6:45PM

How to reach me

- **Office hours** via direct messaging (Discord), video conferencing (Zoom)
- **Announcements** in `#cs450-news` (Discord)
- **Q&A** in `#cs450` (Discord)

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How we are doing remote teaching

- **Open door policy, via Discord.**
  - Message me at any time with your questions.
  - Channel questions answered first, direct-messages answered second.
  - I reply as soon as possible, during office hours in the latest.
- **Homework assignments** we use a grading server (Gradescope)

Course webpage

cogumbreiro.github.io/teaching/cs450/f22/

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Syllabus

cogumbreiro.github.io/teaching/cs450/f22/syllabus.pdf

A	*A-*	*B+*	*B*	*B-*	*C+*	*C*	*C-*	*D+*	*D*	*D-*	*F*
100..95	94..90	89..85	84..80	79..75	74..70	69..65	59..55	54..50	49..45	44..40	39..0

- Course divided into 8 modules
- 1 homework assignment per module
- Final grade: 90% homework + 10% participation

- ***Homework grade:*** average of **8 assignments** (possibly weighted)
- ***Participation grade:*** **in-class quizzes**, attendance classroom/online, participation in forum
- To get `D-` (`C-`) you need to have at least 7 assignments with `D-` (`C-`)

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Academic dishonesty

Plagiarism in University

Copying code from others is wrong because:

- you do not learn
- you risk being expelled
- you are risking the other person being expelled
- you risk not completing your degree
- you risk being put on a list of cheaters (other universities may reject your application)

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Plagiarism in the Industry

Is wrong, because:

- it is illegal
- you risk being dismissed from employment
- you risk being sued

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Copying code (when it is right)

- software licenses define clear rules on
  how you can copy, use, and change other people's code
- open source promotes sharing of code
  - attribution is important (unless public domain)
  - good way to land on a job

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Plagiarism in CS 450

- student's responsibility to learn the Student's code of conduct
- we use plagiarism detection (renaming functions is not enough)
- we compare against solutions from past years (and instructor)
- be careful when working with others, any sharing code may trigger
- the plagiarism detection tool can detect code sharing among students

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Plagiarism in CS 450

Zero Tolerance

- statistically, there will be plagiarism this semester
- if I contact you regarding plagiarism, there will be zero tolerance:
  - You will get an **F** in this course
  - You will be reported to the university

If you need more time to complete an assignment, **ASK**

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Course requirements

Course requirements

Checklist

- Install Racket 7.3: `racket-lang.org`
- Sign in on GitLab (invitation by email)
- Sign in on Discord, say "Hi" in `#cs450-lounge` (invitation link in the GitLab page)
- Sign in on Gradescope, upload the template `hw1.rkt` (invitation by email)

Heads up

- Please, **register using your UMB email address**, otherwise you won't be able to submit your first homework.
- The deadline of homework assignment `n` is last class of module `n` plus 1 week

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Why learn the Structure of Higher Level Languages?

Structure of Higher Level Languages

> I postponed this discussion, because I felt that you are now better suited to understand and relate to the points being made.

- Why learn the fundamental concepts in all programming languages?
- Why learn different languages?
- Why focus on functional programming?
- Why use Racket?

Disclaimer

- Most of these claims are opinions
- These will be mostly informal claims
- We are **not** trying to find the best language (or programming model)

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Overview

- Languages are just tools, learn which language is amenable to what context
- The best programming language does not exist (theoretically most languages are equivalent)
- Different languages have different characteristics that favour different domains: for instance, functional languages being used in Programming Language research, C/Fortran in scientific/high-performance computing
- A programming language is a computing interface: it is crucial to understand its meaning
- The importance of first-class functions and avoiding mutation

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Semantics and idioms

Why should we care about language semantics?

- **A language is a *computing user interface*.**
  We are learning reusable, cross-cutting patterns.
- **The semantics must be *unambiguous* and *precise*.**
  It is not a matter of personal opinion how a conditional expression works. Language features must be described unambiguously to users.
- **The semantics defines a software contract.**
  Is the bug in the client's bug, or is it in our code?
- **Language idioms (patterns) are transferrable knowledge.**
  Understanding idioms (patterns) teaches you something that can be applied across languages and technologies.

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How are all languages similar?

How are all languages the same?

- **Theoretical:** Any input-output behavior implementable in language X is implementable in language Y (Church-Turing thesis), and ***equivalent to the λ-calculus without numbers***
- **Practical:** Reoccurring fundamentals: variables, abstraction, recursive definitions

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How are languages different?

Disclaimer

Languages are not slow/fast

- A language **implementation** is fast/slow, not the language itself
- Certain languages computational models are more amenable to implement efficiently
- Languages are user interfaces of computational models

How different languages behave in different contexts?

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Why is C faster than all other languages?

> Is it because C is "close to the metal?" That is, is C fast because its semantics matches the processor's semantics?

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Why is C faster than all other languages?

> Is it because C is "close to the metal?" That is, is C fast because its semantics matches the processor's semantics? **No!**

- Which processor? How could it match the semantics of all processors?
- Which compiler? The key of C's success lays in having good compilers.
- C compilers are fast because C is **old and its interface remains stable**!
- Popular C compilers are **really** good at optimizing the target language.
- There is a set of good practices to write optimizer-ready C code

Take away

> The facts above make C quite successful in High Performance Computing (large scale scientific codes).

Source: ***C Is Not a Low-level Language: Your computer is not a fast PDP-11***. David Chisnall. ACM Queue vol. 16, no. 2. 2018

Why is Python slow multithreading?

- CPython (the main implementation of Python) is conditioned by the GIL (the Global Interpreter Lock) which effectively serializes parallel execution
- To parallelize code we must run multiple processes, where shared memory is especially slow, which, in turn, slows down compute-bound programs

Take away

> Avoid running compute-bound parallel codes in Python. Maybe choose C?

Source: Global Interpreter Lock. Python Wiki. Last edit in 2017, accessed in 2019.

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Constraint language programming

> We solve the equation `SEND+MORE=MONEY` where each letter represents a digit in Prolog using a constraint language programming module:

```
sendmore(Digits) :-           % Source: https://en.wikipedia.org/wiki/Constraint_programming
   Digits = [S,E,N,D,M,O,R,E],     % Create variables
   Digits ins 0..9,                % Associate domains to variables
   S #\= 0,                        % Constraint: S must be different from 0
   M #\= 0,
   all_different(Digits),          % all the elements must take different values
               1000*S + 100*E + 10*N + D     % Other constraints
             + 1000*M + 100*O + 10*R + E
   #= 10000*M + 1000*O + 100*N + 10*E + Y,
   label(Digits).                  % Start the search
```

Take away

> Some problems are more amenable to certain programming languages.

How are languages different?

1. **The implementation matters:** A language implementation may be conditioned (faster/slower) in certain contexts
2. **The model matters:** Certain problems are simpler/more efficient to write in specific languages
3. **The domain matters:** A technology your business needs may only be available in some language (say TensorFlow in Python)

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Why learn different languages?

> Learn at least one new language every year.

Source: ***The Pragmatic Programmer.*** Andrew Hunt and David Thomas. 1999.

Why should you care

- Deeper understanding of the differences and the similarities between languages
- Learn different approaches to the same problems
- More job opportunities
- Better technology choices (some technologies are only available in specific languages)

Why functional programming?

What is functional programming?

- Mutation is discouraged
- Higher-order functions serve as a generalization device

Why should we care?

- These features help designing correct, elegant, and efficient software
- Functional programming languages are heavily favoured by PL researchers, which means they serve as **a test bed for PL design**. Functional programming is close(r) to math formalism, thus implementation is usually simpler in functional programming languages.
- **Functional programming is trendy!** C++/Java/C#/Python/Javascript are all incorporating functional programming idioms.

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Why should we discourage mutation?

- Simpler to reason about: no surprises passing a data-structure to functions/objects
- Concurrency-ready: read-only means no race conditions (and no locks), which leads to simpler, faster code

Who is using it?

- `immutable.js` for JavaScript by Facebook
- `vavr`, `PCollections`, the Scala runtime, and the Closure runtime for Java
- `immer` for C++
- immutable collections for .NET

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Why should we use higher-order functions?

- Simpler interface than objects (which method? which order?)
- Can be combined effectively (frameworks on combining functions)

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A researcher's Petri Dish

Most programming languages features started out in functional programming languages.

- Garbage collection (LISP, 1959)
- Generics (Hindley-Milner-Damas type system 1969/1978, implemented in ML in ~1977)
- Higher-order functions (`lambda` expressions in C++, C#, Java, Python) introduced in LISP (1959) and in ISWIM (1966)
- Type inference, ***e.g.***, `auto` in C++, `var` in C# (Hindley-Milner-Damas)
- Algebraic-data types and pattern matching (1970s in Hope)
- Recursion

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A new wave of languages

Many new interesting programming languages

- Swift: next-generation programming language for Apple systems
- Rust: functional programming meets system programming
- F#: an ML derivate for the .NET ecosystem
- Elixir: highly-available distributed system
- Clojure: a LISP-influenced language for the JVM and the web

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How are we using functional programming ?

- **OCaml:** web development (Facebook/Meta), distributed systems (Docker), finance (Tezos, Jane Street, Bloomberg, Aesthetic Integration), hardware virtualization (Citrix)
- **Haskell:** verification (Facebook), distributed systems (Google), compilers (Intel), distributed systems (Microsoft)
- **Erlang:** communication (WhatsApp), ads (AddRoll), web backend (Bet365), finance (Goldman Sachs)
- **Elixir:** spam prevention (Pinterest), micro services (Lonely Planet)
- **F#:** data analysis (Kaggle), trading (Credit Suisse), gaming backend (GameSys)
- **Racket** game scripting (Naughty Dog), image processing (YouPatch)
- **Scala** middleware (Twitter), database (Netflix), microservices (Tumblr), web (The Guardian)

Honorable mentions

- ReasonML, Elm, PureScript, ClojureScript

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Course overview

This course is **NOT**...

- **on algorithms**
  For a nice free book read Algorithms by Jeff Erickson.
- **an introduction on programming and computing**
  For a nice free book read How to design programs by Matthias Felleisen, Robert Bruce Findler, Matthew Flatt, Shriram Krishnamurthi
- **on programming with Racket**
  For a nice free book read The Racket Guide by Matthew Flatt, Robert Bruce Findler, and PLT

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This course is...

- **on designing programming language features**
  We will focus mainly on functional and object-oriented programming.
- **on semi-formal specification**
  We will drive our course with precise mathematical notations and tests.
- **on programming patterns**
  We will characterize patterns and study abstractions of these patterns.
- **on purely functional programming**
  We will approach programming without using assignment (mutation).

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Today we will learn

- a formalism to describe a programming language (Racket)
- the semantics of a programming language

How we will learn it

> We introduce one language feature at a time

1. **Syntax:** We formalize each language feature (What)
2. **Example:** We illustrate a feature with an example
3. **Semantics:** We introduce how each language feature works (How)

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Semantics

- Abstract ***Syntax:*** how we write something. Example, which characters/string we use write a keyword, or a number.
- ***Semantics:*** what that something does/means (evaluation here means as the program runs)

> In this class, we focus on the **semantics** of programming languages. We define the semantics of some programming language features.

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1. We shall ***not*** print to output!

Instead, we will use ***assertions***.

2. We shall ***not*** mutate variables!

Instead, we will use ***persistent data structures***.

3. We shall ***not*** use loops!

Instead, we will use ***recursion***.

Your first program

Program

In Racket, **everything evaluates down to or is a value**. A Racket program consists of a preamble followed by zero or more expressions:

```
program = #lang racket expression*
```

1. Racket has no end-of-sentence delimiters (contrary to, say, C-like languages which use semi-colons)
2. Racket evaluates each expression from top-to-bottom, left-to-right

> For space-constraint reasons, code listings might omit the preamble.

Language specification

- `Grayed out text` represents the concrete syntax
- *`Italic text`* represents a meta-variable

Expressions

Expressions can be values, among other things

```
expression = value | ...
```

Values

- **Numbers**
- Void
- Booleans
- Lists
- ...

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Numbers

Numbers

> ***All numbers are complex numbers.*** Some of them are real numbers, and all of the real numbers that can be represented are also rational numbers, except for `+inf.0` (positive infinity), `+inf.f` (single-precision variant), `-inf.0` (negative infinity), `-inf.f` (single-precision variant), `+nan.0` (not-a-number), and `+nan.f` (single-precision variant). Among the rational numbers, some are integers, because round applied to the number produces the same number.

Source: Racket Manual, Section 4.2

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Hello, Numbers!

Your first Racket program

```
#lang racket
10      ; A positive number
+10     ; The plus sign is optional
-10     ; A negative number
0+1i    ; A complex number
1/3     ; A rational number
0.33    ; A floating-point number
```

```
$ racket nums.rkt
10
10
-10
0+1i
1/3
0.33
```

> **Note:** a semi-colon (;) initiates a comment section, which is ignored in Racket. A semi-colon is **not** a end-of-line marker, like in C-like languages.

Expressions are separated by white-space

These two programs are equal:

```
#lang racket
10
+10
-10
0+1i
1/3
0.33
```

```
#lang racket
10 +10 -10     0+1i 1/3 0.33
```

Caveats: `-1` is different than `- 1` (notice the white space in between both characters). The former is the negative one, the latter is the expression `-` and the value `1`. Similarly, `1/3` is a single rational number, whereas `1 / 3` are three expressions.

Function calls

Function call

Delimited by parenthesis and its constituents are separated by white-space characters. The first expression must evaluate to a function, the remaining expressions are the arguments. Each expression is evaluated to a value from left-to-right before applying the function.

```
expression = value | variable | function-call | ...
function-call = ( expression-func expression-arg* )
```

For instance, function call `(expt 2 3)`, for exponentiation, returns `2` raised to the power of `3`. Function `sin` computes the sine function of its sole argument.

```
#lang racket
(expt 2 3)
(sin (expt 2 3))
```

```
$ racket nums-func.rkt
8
0.9893582466233818
```

Note: Function calls can be compounded, as the parameters of a function are arguments too.

No infix notation in Racket

There is **NO INFIX NOTATION** for arithmetic operations (unlike most languages).

The usual arithmetic operations are all just variables: addition `+`, subtraction `-`, multiplication `*`, division `/`.

Example:

```
( * 3.14159 ( * 10 10))
| | |       | | |  |-> Number
| | |       | | |-> Number
| | |       | |-> Variable
| | |       |-> Function call
| | |-> Number
| |-> Variable
|-> Function call
```

> **Note:** In Racket parenthesis represent function application. Contrasted with most C-like languages where parenthesis in expressions are optional and only there to help the reader.

Evaluating a function call

Evaluation works from left-to-right from top-to-bottom

```
#racket lang
; Version 1:
(* 3.14159 (* 10 10))
; Version 2:
(* 3.14159 100)
;          ^^^- Evaluated (* 10 10)
; Version 3:
314.159
;^^^^^^- Evaluated (* 3.14159 * 100)
```

Evaluating a function call

Evaluating a function call

Evaluation works from left-to-right from top-to-bottom

```
#racket lang
; Version 1:
(* 3.14159 (* 10 10))
; Version 2:
(* 3.14159 100)
;          ^^^- Evaluated (* 10 10)
; Version 3:
314.159
;^^^^^^- Evaluated (* 3.14159 * 100)
```

Arithmetic expressions example

$$\Big((11 \cdot 15) + (14 + 4)\Big) + \Big(\frac{3}{9} - (14 \cdot 3)\Big)$$

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Arithmetic expressions example

$$((11 \cdot 15) + (14 + 4)) + \left(\frac{3}{9} - (14 \cdot 3)\right)$$

```
(+
  (+
    (* 11 15)
    (+ 14 4))
  (-
    (/ 3  9)
    (* 14 3)))
```

A longer example

```
(+
  (+
    (* 11 15)
    (+ 14 4))
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  (+
    165
    (+ 14 4))
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  (+
    165
    18)
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  183
  (-
    (/ 3  9)
    (* 14 3)))
```

A longer example

```
(+
  (+
    (* 11 15)
    (+ 14 4))
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  (+
    165
    (+ 14 4))
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  (+
    165
    18)
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  183
  (-
    (/ 3  9)
    (* 14 3)))
```

```
(+
  183
  (-
    1/3
    (* 14 3)))
```

```
(+
  183
  (-
    1/3
    42))
```

```
(+
  183
  -125/3)
```

```
424/3
```

Interpreting an error in Racket

What would happen if we call a function using the infix notation?

```
(3 / 9)
```

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Interpreting an error in Racket

What would happen if we call a function using the infix notation?

```
(3 / 9)
```

```
; application: not a procedure;
;  expected a procedure that can be applied to arguments
;   given: 3
; [,bt for context]
```

Interpreting an error in Racket

What would happen if we call a function using the infix notation?

```
(3 / 9)
```

```
; application: not a procedure;
;  expected a procedure that can be applied to arguments
;   given: 3
; [,bt for context]
```

Line 1

The ***subject*** is `application`. Application is short for function application, aka **calling a function**.

The ***symptom*** is `not a procedure`. Something that should be a procedure is not. Recall, procedure = **function**.

Interpreting an error in Racket

What would happen if we call a function using the infix notation?

```
(3 / 9)
```

```
; application: not a procedure;
;  expected a procedure that can be applied to arguments
;   given: 3
; [,bt for context]
```

Line 1

The ***subject*** is `application`. Application is short for function application, aka **calling a function**.

The ***symptom*** is `not a procedure`. Something that should be a procedure is not. Recall, procedure = **function**.

Line 2

Calling a function requires a function, but we provided something else.

Interpreting an error in Racket

What would happen if we call a function using the infix notation?

```
(3 / 9)
```

```
; application: not a procedure;
;  expected a procedure that can be applied to arguments
;   given: 3
; [,bt for context]
```

Line 1

The ***subject*** is `application`. Application is short for function application, aka **calling a function**.

The ***symptom*** is `not a procedure`. Something that should be a procedure is not. Recall, procedure = **function**.

Line 2

Calling a function requires a function, but we provided something else.

Line 3

We see what was given instead (number 3, rather than a function).

Is this example a legal Racket program?

```
#lang racket
sin
```

Is this example a legal Racket program?

```
#lang racket
sin
```

Yes! `sin` is a variable, so a valid expression. Hence, Racket just prints what is in variable `sin`.

```
$ racket sin.rkt
#<procedure:sin>
```

> **Note:** In Racket lingo the word ***procedure*** is a synonym for function.

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Racket specification

```
program = #lang racket expression*
expression = value | variable | function-call | ...
value = number | ...
function-call = ( expression+ )
```