

CS420

Introduction to the Theory of Computation

Lecture 9: Power, Kleene star, equivalence

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Today we will learn...

- Void
- All
- Power
- Kleene star
- Language equivalence

The void language

Void

■ The language that rejects all strings.

Void

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Definition $\text{Void } w := \text{False}$.

Correction properties

1. Show every word is rejected by Void

The all language

All

Language that accepts all strings

All

Language that accepts all strings

Definition All (w :word) := True.

Correction properties

1. Show that any word is accepted by All.

Exercises

Coq	Notation	Math
Nil		$\{\epsilon\}$
Char c	c	$\{c\}$
Union L1 L2	L1 U L2	$L_1 \cup L_2$
App L1 L2	L1 >> L2	$L_1 \cdot L_2$
Void		\emptyset
All		Σ^*

Solve the following exercises

1. $L_1 \cup \{\epsilon\} =$

- $L_1 = \{[0], [1], [2]\}$
- $L_2 = \{[3], [4]\}$

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 $\{[0, 3], [0, 4], [1, 3], [1, 4], [2, 4], [2, 5]\}$
4. $L_2 \cdot \{\epsilon\} =$

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7. $L_2 \cdot \emptyset = \emptyset$

The power operator for languages

The power operator for languages

- $L^{n+1} = L \cdot L^n$
- $L^0 = \{\epsilon\}$

Example

- $L = \{[0], [1], [2]\}$
- $L^0 = \{\epsilon\}$
- $L^1 = L \cdot \{\epsilon\} = L$
- $L^2 = L \cdot L = \{[0, 0], [0, 1], [0, 2], [1, 0], [1, 1], [1, 2], [2, 0], [2, 1], [2, 2]\}$

Implementing power

```

Inductive Pow (L:language) : nat → word → Prop :=
| pow_nil:
  Pow L 0 nil
| pow_cons:
  forall n w1 w2 w3,
  In w2 (Pow L n) →
  In w1 L →
  w3 = w1 ++ w2 →
  Pow L (S n) w3.

```

Rule pow_nil:

$$\epsilon \in L^0$$

Rule pow_cons:

$$\frac{w_2 \in L^n \quad w_1 \in L}{w_1 \cdot w_2 \in L^{n+1}}$$

Rules in the form of:

$$\frac{P_1 \quad P_2 \quad P_3}{Q}$$

Are read as: If P_1 **and** P_2 **and** P_3 all hold, **then** we have Q .

Exercise

```
Require Import Coq.Strings.Ascii.  
Require Import Coq.Lists.List.  
From Turing Require Import Lang.  
From Turing Require Import Util.  
Import Lang.Examples.  
Import LangNotations.  
Import ListNotations.  
Open Scope lang_scope.  
Open Scope char_scope.
```

Lemma in_aaa:

```
  In ["a"; "a"; "a"] (Pow "a" 3).
```

Proof.

Qed.

Lemma pow_char_in_inv:

```
  forall c n w,  
  In w (Pow (Char c) n) →  
  w = Util.pow1 c n.
```

Proof.

Qed.

Kleene operator

Kleene operator

$$L^* = L^0 \cup L^1 \cup L^2 \cup L^3 \cup \dots$$

Inductive definition

$$\frac{w \in L^n}{w \in L^*}$$

Wait, what is n ?

Any n will do. If you can build a proof object such that $w \in L^n$, then $w \in L^*$.

Does this mean that there is only one n ? Say, $L^* = L^{1000}$?

NO it does not. Each word membership will have its possibly distinct n .

Example: $L = [a]$, we have that $\epsilon \in L^0$ and that $[a, a] \in L^2$, thus $\epsilon \in L^*$ and $[a, a] \in L^*$.

Exercise

Lemma `in_aaa_2`:

In `["a"; "a"; "a"]` (Star `"a"`).

Proof.

Language Equivalence

Language equivalence (equality)

- Mathematically, we write $L_1 = L_2$ to mean that two languages are equal.
- How do you prove language equality?

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- Mathematically, we write $L_1 = L_2$ to mean that two languages are equal.
- How do you prove language equality?
- You have to show that all words in L_1 are also in L_2 and vice-versa.

Language equivalence in Coq

Definition `Equiv (L1 L2:language) := forall w, L1 w ↔ L2 w.`

Show that `Vowel` is equivalent to previous example

Lemma `vowel_eq:`

`Vowel == (Char "a" U Char "e" U Char "i" U Char "o" U Char "u").`

Proof.

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Lemma `vowel_eq:`

`Vowel == (Char "a" U Char "e" U Char "i" U Char "o" U Char "u").`

Proof.

`apply vowel_iff.`

Qed.

Exercise

Show that `Void` is a neutral element in union.

```
Lemma union_l_void:  
  forall L,  
    L U Void == L.
```

Exercise

Show that Void is a neutral element in union.

Lemma union_l_void:

```
forall L,
  L U Void == L.
```

Proof.

```
split; intros.
- destruct H. {
  assumption.
}
apply not_in_void in H.
contradiction.
- left.
  assumption.
```

Qed.

Exercise

Show that `Void` is an absorbing element in concatenation.

```
Lemma app_l_void:  
  forall L,  
    L >> Void = Void.
```

Exercise

Show that Void is an absorbing element in concatenation.

```
Lemma app_l_void:  
  forall L,  
    L >> Void = Void.
```

Proof.

```
unfold App; split; intros.  
- destruct H as (w1, (w2, (Ha, (Hb, Hc)))).  
  subst.  
  apply not_in_void in Hc.  
  contradiction.  
- apply not_in_void in H.  
  contradiction.
```

Qed.

Exercise

■ A language that accepts any words that consists of two vowels

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Definition `TwoVowels := Vowel >> Vowel.`

Show that `["a"; "e"]` is in `TwoVowels`

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Goal In `["a"; "e"] (Vowel >> Vowel).`

Proof.

Exercise

A language that accepts any words that consists of two vowels

Definition `TwoVowels := Vowel >> Vowel.`

Show that `["a"; "e"]` is in `TwoVowels`

Goal In `["a"; "e"] (Vowel >> Vowel).`

Proof.

```

unfold App.
exists ["a"], ["e"]. (* Existential in the goal *)
split. { reflexivity. }
split. { left. reflexivity. }
right. left. reflexivity.

```

Qed.

Exercise

What words are accepted by L2?

Definition $L2 := All \gg Char \text{"a"}$.

Exercise

■ Rewrite `Vowels` to use only language operators.

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```
Definition Vowels2 := Char "a" U Char "e" U Char "i" U Char "o" U Char "u".
```

Exercise

Lemma `vowel_length`:

`forall` w ,

Vowel $w \rightarrow$

`length` $w = 1$.

Exercise

Lemma `vowel_length`:

```
forall w,  
Vowel w →  
length w = 1.
```

Proof.

```
intros.  
destruct H as [H|[H|[H|[H|H]]]]; subst; reflexivity.
```

Qed.

Exercise

Goal for all w , $(\text{Vowel} \gg \text{Vowel}) w \rightarrow \text{length } w = 2$.

Exercise

Goal forall w, (Vowel >> Vowel) w → length w = 2.

Proof.

intros.

unfold App in *.

destruct H as (w1, (w2, (Ha, (Hb, Hc)))). (* Existential in hypothesis *)

subst. apply vowel_length in Hb. apply vowel_length in Hc.

SearchAbout (length(_ ++ _)). (* Search for lemmas *)

rewrite app_length. rewrite Hb. rewrite Hc. reflexivity.

Qed.

Exercise

Show that all strings are rejected by `Void`.

Exercise

Show that all strings are rejected by Void.

Lemma not_in_void:

```
forall w,  
  ~ In w Void.
```

Proof.

```
intros.  
intros N.  
inversion N.
```

Qed.