

CS420

Introduction to the Theory of Computation

Lecture 11: Regular expressions

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Today we will learn...

- Language equivalence theorems
- Summary of Lang. Examples
- Regular expressions

Language equivalence

- $L_1 \cup (L_2 \cup L_3) = (L_1 \cup L_2) \cup L_3$

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- $(L \cdot \emptyset) = (\emptyset \cdot L) = \emptyset$

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- $L \cup \mathbf{All} = \mathbf{All} \cup L = \mathbf{All} = \Sigma^*$

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- $(L \cdot \emptyset) = (\emptyset \cdot L) = \emptyset$
- $(L \cdot \epsilon) = (\epsilon \cdot L) = (L \cup \emptyset) = (\emptyset \cup L) = (L \cup L) = L$
- $L \cup \mathbf{A11} = \mathbf{A11} \cup L = \mathbf{A11} = \Sigma^*$
- $L^* \cdot L^* = (L^*)^* = L^*$

Language equivalence

- $L_1 \cup (L_2 \cup L_3) = (L_1 \cup L_2) \cup L_3$
- $(L_1 \cup L_2) = (L_2 \cup L_1)$
- $(L \cdot \emptyset) = (\emptyset \cdot L) = \emptyset$
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- $L \cup \mathbf{All} = \mathbf{All} \cup L = \mathbf{All} = \Sigma^*$
- $L^* \cdot L^* = (L^*)^* = L^*$
- $\epsilon^* = \emptyset^* = \epsilon$

Language equivalence

- $L_1 \cup (L_2 \cup L_3) = (L_1 \cup L_2) \cup L_3$
- $(L_1 \cup L_2) = (L_2 \cup L_1)$
- $(L \cdot \emptyset) = (\emptyset \cdot L) = \emptyset$
- $(L \cdot \epsilon) = (\epsilon \cdot L) = (L \cup \emptyset) = (\emptyset \cup L) = (L \cup L) = L$
- $L \cup \mathbf{All} = \mathbf{All} \cup L = \mathbf{All} = \Sigma^*$
- $L^* \cdot L^* = (L^*)^* = L^*$
- $\epsilon^* = \emptyset^* = \epsilon$
- $(L_1 \cdot L_3) \cup (L_2 \cdot L_3) = ((L_1 \cup L_2) \cdot L_3)$

Lang. v examples

(Live coding...)

Regular expressions

Use Case 2

Regular Expressions: Input validation

Use Case 2

Regular Expressions: Input validation

HTML includes regular expressions to perform client-side form validation.

```
<input id="uname" name="uname" type="text"
  pattern="_([a-z]|[A-Z]|[0-9])+" minlength="4" maxlength="10">
```

- `_[a-zA-Z0-9]+`
- `[a-zA-Z0-9]` means any character between a and z, or between A and Z, or between 0 and 9
- `R+` means repeat R one or more times
- In this case, the username must start with an underscore `_`, and have one or more letters/numbers
- `minlength` and `maxlength` further restrict the string's length

Use Case 3

Regular Expressions: Text manipulation

Use Case 3

Regular Expressions: Text manipulation

Programming languages include regular expressions for fast and powerful text manipulation.

Example (JS)

```
let txt1 = "Hello World!";  
let txt2 = txt1.replace(/[a-zA-Z]+/, "Bye"); // Replaces the first word by "Bye"  
console.log(txt2);  
// Bye World!
```


Regular expressions

A theoretical motivation

What languages can we specify with the following operators?

- Void
- Nil
- Char
- App
- Union
- Star

Regular expressions

A theoretical motivation

What languages can we specify with the following operators?

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Idea: specify a datatype that represents all possible expressions

Regular expressions

We define regular expression R as either:

Notation	Meaning
\emptyset	Rejects all words
ϵ	Only accepts the empty string
c	Only accepts a string with a single character c
$R_1 R_2$	Accepts a word from R_1 concatenated with a word from R_2
$R_1 R_2$	Accepts a word from R_1 or a word from R_2
R^*	Accepts 0 or more copies of words from R

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`
Accept

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`
Accept
- **Input:** `[b,a,r]`

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`
Accept
- **Input:** `[b,a,r]`
Accept
- **Input:** `[f,o]`

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`
Accept
- **Input:** `[b,a,r]`
Accept
- **Input:** `[f,o]`
Reject
- **Input:** `[]`

Exercise

Expression: `foo || bar`

■ Is the following input accepted or rejected?

- **Input:** `[f,o,o]`
Accept
- **Input:** `[b,a,r]`
Accept
- **Input:** `[f,o]`
Reject
- **Input:** `[]`
Reject

Exercise

Expression: $(\mathbf{foo||bar})^*$

- **Input:** [f,o,o]

Exercise

Expression: $(\mathbf{foo}||\mathbf{bar})^*$

- **Input:** [f,o,o]
Accept

Exercise

Expression: $(\mathbf{foo}||\mathbf{bar})^*$

- **Input:** [f,o,o]
Accept
- **Input:** [b,a,r]

Exercise

Expression: $(\text{foo}||\text{bar})^*$

- **Input:** [f,o,o]
Accept
- **Input:** [b,a,r]
Accept
- **Input:** [f,o]

Exercise

Expression: $(\text{foo}||\text{bar})^*$

- **Input:** [f,o,o]
Accept
- **Input:** [b,a,r]
Accept
- **Input:** [f,o]
Reject
- **Input:** []

Exercise

Expression: $(\text{foo}||\text{bar})^*$

- **Input:** [f,o,o]
Accept
- **Input:** [b,a,r]
Accept
- **Input:** [f,o]
Reject
- **Input:** []
Accept
- **Input:** [f,o,o,b,a,r]

Exercise

Expression: $(\text{foo}||\text{bar})^*$

- **Input:** [f,o,o]
Accept
- **Input:** [b,a,r]
Accept
- **Input:** [f,o]
Reject
- **Input:** []
Accept
- **Input:** [f,o,o,b,a,r]
Accept

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b]

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b] **REJECT**
- [b,c,a]

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT

Exercise

Expression

$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$

- [b] REJECT
- [b,c,a] REJECT
- [c]

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b]

Exercise

Expression

$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- []

Exercise

Expression

$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] REJECT
- [a,a]

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] REJECT
- [a,a] ACCEPT
- [a]

Exercise

Expression

$$(\emptyset \cdot c) \parallel aa \parallel a \cdot \epsilon$$

- [b] REJECT
- [b,c,a] REJECT
- [c] REJECT
- [a,b] REJECT
- [] REJECT
- [a,a] ACCEPT
- [a] ACCEPT

Examples

Source: regexlib.com

This expression matches a hyphen separated US phone number, of the form ANN-NNN-NNNN, where A is between 2 and 9 and N is between 0 and 9.

- `[2-9][0-9]{2}-[0-9]{3}-[0-9]{4}`

Examples

Source: regexlib.com

This expression matches a hyphen separated US phone number, of the form ANN-NNN-NNNN, where A is between 2 and 9 and N is between 0 and 9.

- `[2-9][0-9]{2}-[0-9]{3}-[0-9]{4}`

Breaking it down:

- `[2-9]` corresponds to 2 || 3 || 4 || 5 || 7 || 8 || 9
- `[0-9]{2}` corresponds to the power of 2, thus pattern `[0-9][0-9]`
- -
- `[0-9]{3}`
- -
- `[0-9]{4}`