

CS420

Logical Foundations of Computer Science

Lecture 7: Mock mini-test 1

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Today we will learn...

- Existential operator
- Mock Mini-Test 1
- Formal language
- Language operators
- Language equivalence

From proposition to proof state

```
Goal forall (a b c:nat), a = b → b = c.
```

```
Proof.
```

```
  intros.
```

What is the expected proof state?

From proposition to proof state

Goal forall (a b c:nat), a = b → b = c.

Proof.

intros.

What is the expected proof state?

Solution

```
1 subgoal
a, b, c : nat
H : a = b
----- (1/1)
b = c
```

- Each parameter of a theorem is an **assumption**
- Each **variable** in the forall is one parameter becomes an assumption
- Each **pre-condition** of an implication becomes an assumption
- Variables and pre-conditions are parameters

You can name assumptions in a forall

```
Goal forall (a b c:nat) (eq_a_b: a = b),
```

```
  b = c.
```

```
Proof.
```

```
  intros.
```

What is the expected proof state?

You can name assumptions in a forall

```
Goal forall (a b c:nat) (eq_a_b: a = b),
```

```
  b = c.
```

```
Proof.
```

```
  intros.
```

What is the expected proof state?

Solution

```
1 subgoal
```

```
a, b, c : nat
```

```
eq_a_b : a = b
```

```
----- (1/1)
```

```
b = c
```

- Implications are just **anonymous** parameters (name will be generated automatically)
- Think assert $(x = y)$ versus assert $(\text{Ha}: x = y)$

From proof state to proposition:

What is the lemma that originates the following proof state?

```

a, b, c: nat
P, Q: Prop
H: P → a = b
H0: Q ∨ P
H1: b = c
----- (1/1)
a = c
  
```

From proof state to proposition:

What is the lemma that originates the following proof state?

```

a, b, c: nat
P, Q: Prop
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H0: Q ∨ P
H1: b = c
----- (1/1)
a = c

```

Solution 1:

```
Goal forall (a b c: nat) (P Q: Prop) (H: P → a = b) (H0: Q ∨ P) (H1: b = c), a = c.
```

Solution 2:

```
Goal forall (a b c: nat) (P Q: Prop), (P → a = b) → (Q ∨ P) → (b = c) → a = c.
```


Existential quantification

$$\exists x.P$$

Existential quantification

```
Inductive ex (A : Type) (P : A → Prop) : Prop :=
  | ex_intro : forall (x : A) (h : P x), ex P.
```

Notation:

```
exists x:A, P x
```

- To conclude a goal `exists x:A, P x` we can use tactics `exist x.` which yields `P x`. Alternatively, we can use `apply ex_intro`.

```
forall n, exists z, z + n = n
```

- To use a hypothesis of type `H:exists x:A, P x`, you can use `destruct H as (x,H)`, or `inversion H`

```
forall n, (exists m, m < n) → n <> 0.
```

Defining arbitrary logical relations

Defining less-than-equal

Inductive definition of \leq

$$\frac{}{n \leq n} \mathbf{le_n}$$

$$\frac{n \leq m}{n \leq S m} \mathbf{le_S}$$

```

Inductive le : nat → nat → Prop :=
| le_n : forall n:nat,
  n ≤ n
| le_S : forall (n m : nat),
  n ≤ m →
  n ≤ S m.
  
```

- Any pre-condition will appear above the line
- Preconditions are separated by whitespace

How do we know that less-than-equal was defined correctly?



How do we know that less-than-equal was defined correctly?

With theorems!

```

(* Simple tests *)
Goal 1 ≤ 1. Proof. Admitted.
Goal 1 ≤ 10. Proof. Admitted.
(* More interesting properties *)
Theorem le_is_reflexive: forall x,
  x ≤ x.
Proof. Admitted. (* Proved in class *)
Theorem le_is_anti_symmetric: forall x y,
  x ≤ y →
  y ≤ x →
  x = y.
Proof. Admitted. (* Proved in class *)
Theorem le_is_transitive: forall x y z,
  x ≤ y →
  y ≤ z →
  x ≤ z.
Proof. Admitted.

```

Mock Mini-Test 1

Q1.1

All functions defined in Coq via `Fixpoint` must terminate on all inputs.

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Solution: True

All functions must terminate.

Q1.2

If $S(n + m) = n + S m$ is the goal in the current proof state, then `{reflexivity}` will solve the goal.

Q1.2

If $S (n + m) = n + S m$ is the goal in the current proof state, then `{reflexivity}` will solve the goal.

Solution: False

Goal

```
forall n m,
  S (n + m) = n + S m.
```

Proof.

```
intros.
```

```
Fail reflexivity.
```

Abort.

Q1.3

A **polymorphic** type is one that is parameterized by a type argument by using the universal quantifier `forall`. For instance: `forall (X:Type), list X → list X` is a polymorphic type.

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Solution: True

Q1.4

If E has type `beq_nat m n = true`, then E also has type `m = n`.

Q1.4

If E has type `beq_nat m n = true`, then E also has type `m = n`.

Solution: False

Goal

```
forall n m (E:Nat.eqb n m = true),
  m = n.
```

Proof.

```
intros.
```

```
Fail apply E.
```

Abort.

Q1.5

The proposition $\forall n, S\ n \leftrightarrow n$ is provable in Coq.

Q1.5

The proposition $\text{forall } n, S\ n \lt;> n$ is provable in Coq.

Solution: True

```

Goal
  forall n, S n <> n
.
Proof.
  intros.
  intros N.
  induction n. {
    inversion N.
  }
  inversion N.
  apply IHn.
  assumption.
Qed.

```

Q2.1

What is the type of the following expression?

Nat.eqb 28

Q2.1

What is the type of the following expression?

```
Nat.eqb 28
```

Answer: `nat → bool`

Q2.2

What is the type of the following expression?

`14 = 68`

Q2.2

What is the type of the following expression?

$14 = 68$

Answer: Prop

Q3.1

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall n, n <> S n
```

Q3.1

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

`forall n, n <> S n`

Answer: induction

Q3.2

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall (n m:nat), n = m \ / n <> m
```


Q3.2

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall (n m:nat), n = m \ / n <> m
```

Answer: BY INDUCTION

Q3.3

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall A B:Type, forall (f g: A → B), f = g → forall x, f x = g x
```

Q3.3

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall A B:Type, forall (f g: A → B), f = g → forall x, f x = g x
```

Answer: EASY

Goal

```
forall A B:Type, forall (f g: A → B), f = g → forall x, f x = g x.
```

Proof.

```
intros.
```

```
rewrite H.
```

```
reflexivity.
```

Qed.

Q3.4

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall P : Prop, P
```

Q3.4

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

```
forall P : Prop, P
```

Answer: NOT PROVABLE

Goal

```
forall P : Prop, P.
```

Proof.

```
intros X.
```

```
Fail apply X.
```

Abort.

Q3.5

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

`forall n, n+5 ≤ n+6`

Q3.5

The proof of this goal is: EASY / BY INDUCTION / NOT PROVABLE

forall n, $n+5 \leq n+6$

Answer: INDUCTION

Q4.1

Prove this goal:

$H : \sim \sim P$

$H0 : P \vee \sim P$

----- (1/1)

P

Q4.1

Prove this goal:

$$H : \sim \sim P$$

$$H0 : P \ \backslash / \ \sim P$$

----- (1/1)
P

```
destruct H0. {
  assumption.
}
apply H in H0.
contradiction.
```

Q4.2

Prove this goal:

$$H : P \rightarrow Q$$
$$H0 : P \vee \sim P$$

----- (1/1)

$$\sim P \vee Q$$

Q4.2

Prove this goal:

$$\begin{array}{l}
 H : P \rightarrow Q \\
 H0 : P \vee \sim P \\
 \hline
 \sim P \vee Q
 \end{array}
 \quad (1/1)$$

```

destruct H0. {
  apply H in H0.
  right.
  assumption.
}
left.
assumption.
  
```

Q4.3

Prove this goal:

$P, Q : \text{Prop}$

$PQ : P \rightarrow Q$

$NQ : \sim Q$

$HP : P$

----- (1/1)

False

Q4.3

Prove this goal:

```
P, Q : Prop
PQ : P → Q
NQ : ~ Q
HP : P
```

----- (1/1)

False

```
apply PQ in HP. contradiction.
```

Q4.4

-----^(1/1)
`forall (A:Type) (l:list A), l = [] → l = []`

Q4.4

```
-----(1/1)  
forall (A:Type) (l:list A), l = [] → l = []
```

```
intros. assumption.
```